

# Special Topics in Many-Valued Logics

Part II - Riesz MV-algebras and their logic

Ioana Leuştean

**Special Topics in Computer Science**

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## The product in Łukasiewicz logic

- internal product:  $([0, 1], \oplus, \cdot, *, 0)$

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A. Di Nola, I.L., 2011: Riesz MV-algebras
- internal product and scalar multiplication:  
 $([0, 1], \oplus, \cdot, *, \{r|r \in [0, 1]\}, 0)$   
S. Lapenta, I.L., 2014: fMV-algebras

## MV-algebras with product

<i>Category</i>	<i>MV-algebras with product</i>	<i>unital <math>\ell</math>-structures</i>	<i>Standard model</i>
<b>MV</b>	MV-algebras	$\ell$ -groups	$\text{HSP}([0, 1])$
<b>RMV</b>	Riesz MV-algebras	Riesz spaces	$\text{HSP}([0, 1])$
<b>PMV</b>	Product MV-algebras	$f$ -rings	$\text{PMV}^+ = \text{ISP}([0, 1])$
<b>fMV</b>	$f$ MV-algebras	$f$ -algebras	$f\text{MV}^+ = \text{ISP}([0, 1])$

## MV-algebras with product

<i>Logic</i>	<i>MV-algebras with product</i>	<i>unital <math>\ell</math>-structures</i>	<i>Standard model</i>
$\mathcal{L}_\infty$	MV-algebras	$\ell$ -groups	$\text{HSP}([0, 1])$
$\mathbb{RL}$	Riesz MV-algebras	Riesz spaces	$\text{HSP}([0, 1])$
$\mathcal{PL}$	Product MV-algebras	$f$ -rings	$\text{PMV}^+ = \text{ISP}([0, 1])$
$f\mathcal{L}$	$f$ MV-algebras	$f$ -algebras	$f\text{MV}^+ = \text{ISP}([0, 1])$

**Riesz MV-algebra**  $(A, \oplus, *, \{r \mid r \in [0, 1]\}, 0_A)$

- ①  $(R, \oplus, *, 0_A)$  is an MV-algebra,
- ②  $r(x \odot y^*) = (rx) \odot (ry)^*$ ,
- ③  $(r \odot q^*) \cdot x = (rx) \odot (qx)^*$ ,
- ④  $r(qx) = (rq)x$ ,
- ⑤  $1x = x$ ,

$r, q \in [0, 1]$  and  $x, y \in A$ , where  $x \odot y = (x^* \oplus y^*)^*$ .

**Riesz MV-algebras and Riesz spaces**

- For any Riesz MV-algebra  $A$  there is a Riesz space with strong unit  $(V, u)$  such that  $A \simeq [0, u]_V$ .
- The category of Riesz MV-algebras is equivalent with the category of Riesz spaces with strong unit.

## The logic $\mathbb{RL}$

The connectives are  $\{\rightarrow, \neg\} \cup \{\sigma_r \mid r \in [0, 1]\}$

The axioms:

- ① the axioms of  $\mathcal{L}$
- ②  $\sigma_r(\varphi \rightarrow \psi) \leftrightarrow (\sigma_r\varphi \rightarrow \sigma_r\psi)$
- ③  $\sigma_{(r \odot q^*)}\varphi \leftrightarrow (\sigma_q\varphi \rightarrow \sigma_r\varphi)$
- ④  $\sigma_r\sigma_q\varphi \leftrightarrow \sigma_{(rq)}\varphi$
- ⑤  $\sigma_1\varphi \leftrightarrow \varphi,$

where  $\varphi, \psi, \chi$  are formulas and  $r, q \in [0, 1]$ .

The deduction rule is *modus ponens*.

### Theorem

The Lindenbaum-Tarski algebra is a Riesz MV-algebra.



## The logic $\mathbb{RL}$

### Completeness. TFAE:

- $\varphi$  is provable,
- $e(\varphi) = 1$  for any  $[0, 1]$ -evaluation.

### Theorem

$$\mathbf{RMV} = \mathbf{HSP}([0, 1])$$

# The logic $\mathbb{RL}$

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## Theorem

$$\mathbf{RMV} = \mathbf{HSP}([0, 1])$$

For  $r \in [0, 1]$  we set  $\mathbf{r} := \neg(\sigma_r(\neg\theta))$  where  $\vdash \theta$

## Pavelka completeness

If  $\varphi$  is a formula of  $\mathbb{RL}$ , we define:

- the *truth degree* of  $\varphi$ , by
$$\|\varphi\| = \min\{e(\varphi) \mid e \text{ is an evaluation}\},$$
- the *provability degree* of  $\varphi$ , by
$$|\varphi| = \max\{r \in [0, 1] \mid \vdash \mathbf{r} \rightarrow \varphi\},$$

then  $|\varphi| = \|\varphi\|$ .

## Normal form theorems

- $\varphi$  formula with  $n$  variables  $\mapsto f_\varphi : [0, 1]^n \rightarrow [0, 1]$

A function  $f : [0, 1]^n \rightarrow [0, 1]$  is a  $\text{PWL}_u(\mathbb{Z})$  function if it is continuous and there is a finite set of affine functions  $p_1, \dots, p_k : \mathbb{R}^n \rightarrow \mathbb{R}$  with integer coefficients such that for any  $(a_1, \dots, a_n) \in [0, 1]^n$  there exists  $i \in \{1, \dots, k\}$  with  $f(a_1, \dots, a_n) = p_i(a_1, \dots, a_n)$ .

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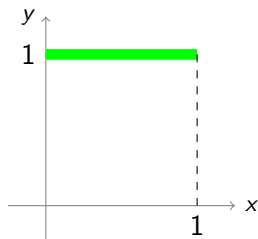
Free MV-algebra  $MV_n \simeq \text{Lind}_{\mathcal{L}, n}$  [R. McNaughton, 1951]

$$\text{Free}_{MV}(n) = \{f_\varphi : [0, 1]^n \rightarrow [0, 1] \mid \varphi \text{ formula of } \mathcal{L}\} = \text{PWL}_u(\mathbb{Z})$$

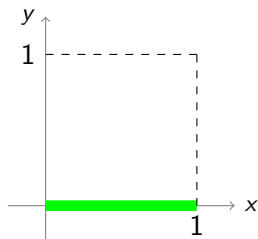
Free Riesz MV-algebra  $RMV_n \simeq \text{Lind}_{\mathbb{R}\mathcal{L}, n}$

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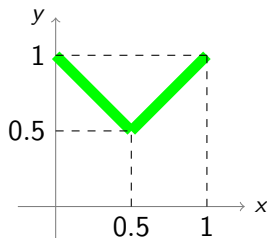
## Elements of $Free_{MV}(1)$



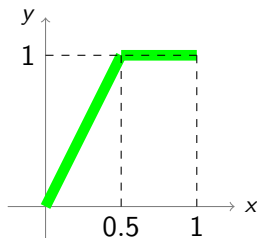
$$f(x) = x \oplus x^*$$



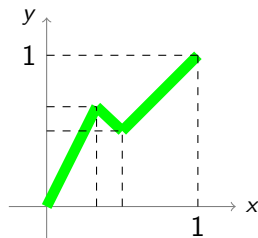
$$f(x) = x \odot x^*$$



$$f(x) = x \vee x^*$$

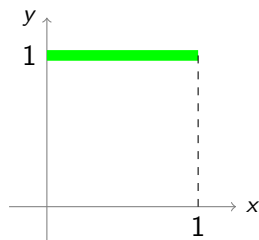


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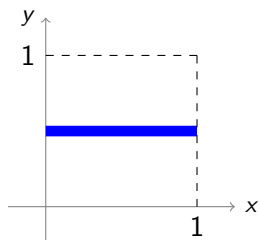


$$f(x) = (x \oplus x) \wedge (x \vee x^*)$$

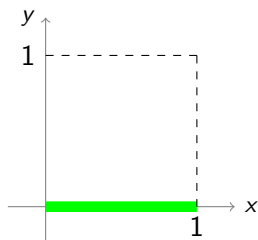
## Elements of $Free_{MV}(1)$ and $Free_{RMV}(1)$



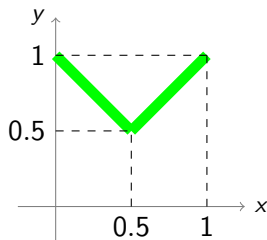
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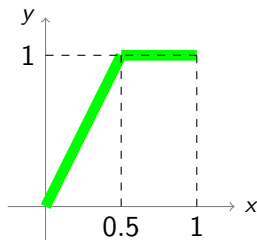
$$\in Free_{RMV}(1)$$



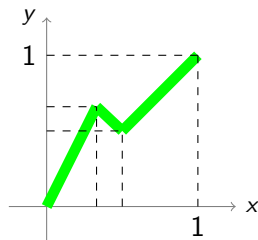
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$$f(x) = (x \oplus x) \wedge (x \vee x^*)$$

- Riesz MV-algebras are a variety,
- Riesz MV-algebras have a simple axiomatization,
- the logic  $\mathbb{R}\mathcal{L}$  is "classically" developed,
- the logic  $\mathbb{R}\mathcal{L}$  has standard completeness and normal-form theorem,
- Riesz MV-algebras categorically equivalent with Riesz spaces with strong unit.

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... most of the standard real function spaces are vector lattices, and in a very natural way.

G. Birkhoff, Lattice Theory, 1973



## Banach lattices: M-spaces, L-space

A Banach lattice  $(R, \|\cdot\|)$  is an

- **M-space:**  $\|x \vee y\| = \|x\| \vee \|y\|$  whenever  $x, y \geq 0$
- **L-space:**  $\|x + y\| = \|x\| + \|y\|$  whenever  $x, y \geq 0$

### Examples

- $C(X, \mathbb{R})$  is an M-space for any compact Hausdorff space  $X$ .
- $L_1(\mu)$  is an L-space for any measure space  $(X, \Omega, \mu)$ .

### Kakutani's theorems

- For any M-space with strong unit  $(M, u)$  there exists a compact Hausdorff space  $X$  such that  $(M, u)$  is isometrically Riesz isomorphic with  $C(X, \mathbb{R})$ .
- For any L-space with strong unit  $(L, u)$  there exists a measure space  $(X, \Omega, \mu)$  such that  $L$  is isometrically Riesz isomorphic with  $L^1(\mu)$ .

## Unit intervals in M-spaces

On any semisimple Riesz MV-algebra  $A$  define

$$\|x\|_u = \inf \{r \in [0, 1] \mid x \leq r \cdot 1\} \text{ for any } x \in A \text{ (unit norm)}$$

Note that  $\|x \vee y\|_u = \|x\|_u \vee \|y\|_u$  for any  $x, y \in A$ .

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### Norm-complete Riesz MV-algebras and M-spaces

Any Riesz MV-algebra which is norm-complete w.r.t  $\|\cdot\|_u$  is isomorphic with the unit interval of an M-space with strong unit.

The category of Riesz MV-algebras which are norm-complete w.r.t  $\|\cdot\|_u$  is equivalent with the category of unital M-spaces with Riesz homomorphisms.

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### Norm-complete Riesz MV-algebras and compact Hausdorff spaces

A Riesz MV-algebra  $A$  is norm-complete w.r.t.  $\|\cdot\|_u$  iff  
 $A \simeq C(\text{Max}(A), [0, 1])$ .

The category of Riesz MV-algebras which are norm-complete w.r.t.  $\|\cdot\|_u$  with Riesz MV-algebra homomorphisms is:

- dually equivalent with the category of compact Hausdorff spaces with continuous functions,
- equivalent with the category of  $C^*$ -algebras.

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I.L.: similar results for L-spaces

$\mathbb{R}\mathcal{L}$  the logic of Riesz MV-algebras

$[\varphi]$  in  $Lind_{\mathbb{R}\mathcal{L},n}$

$$\|[\varphi]\|_u = \sup\{f_\varphi(\mathbf{x}) \mid \mathbf{x} \in [0, 1]^n\}$$

$(Lind_{\mathbb{R}\mathcal{L},n}, \|\cdot\|_u)$  is a normed space

### Norm-completion of the Lindenbaum-Tarski algebra

The norm-completion of the normed space  $(Lind_{\mathbb{R}\mathcal{L},n}, \|\cdot\|_u)$  is isometrically isomorphic with  $(C([0, 1]^n), \|\cdot\|_\infty)$ .

### Approximation of continuous functions

For any continuous function  $f : [0, 1]^n \rightarrow [0, 1]$  there exists a sequence of formulas  $(\varphi_n)_n$  of  $\mathbb{R}\mathcal{L}$  such that  $\lim_n f_{\varphi_n} = f$ .

For  $r \in [0, 1]$  we set  $\mathbf{r} := \neg(\sigma_r(\neg\theta))$  where  $\vdash \theta$

Uniform limit (inspired by [X. Caicedo, LATD'08])

A formula  $\varphi$  is *the uniform limit* of the sequence  $(\varphi_n)_n$  in  $\mathbb{RL}$  if for any  $r < 1$  there is  $k$  such that for any  $n \geq k$ :  $\vdash \mathbf{r} \rightarrow (\varphi \leftrightarrow \varphi_n)$ .

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## The logic $\mathbb{QL}$

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The language is  $\{\rightarrow, \neg\} \cup \{\sigma_r \mid r \in [0, 1] \cap \mathbb{Q}\}$

The axioms are similar with the ones of  $\mathbb{RL}$ .

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## Rational Łukasiewicz logic [B. Gerla, 2001]

The language is  $\{\rightarrow, \neg\} \cup \{\delta_n \mid n \in \mathbb{N}\}$

I. the axioms of Ł:

II. the following formulas are axioms for any  $n \geq 2$ :

$$\varphi \leftrightarrow n\delta_n\varphi$$

$$(\varphi \rightarrow n\psi) \rightarrow (\delta_n\varphi \rightarrow \psi)$$

$\mathbb{QL}$  and *Rational Łukasiewicz logic* are equivalent.

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The axioms are similar with the ones of  $\mathbb{RL}$ .

## DMV-algebras

The models of  $\mathbb{QL}$  are defined as Riesz MV-algebras, but we have only rational scalars.

They are categorically equivalent with:

- divisible MV-algebras,
- $\mathbb{Q}$ -vector lattices with strong unit

$$\mathbf{DMV} = \mathbf{HSP}([0, 1]_{\mathbb{Q}})$$

### Formulas as limits

For any formula  $\varphi$  of  $\mathbb{R}\mathcal{L}$  there exists a sequence of formulas  $(\varphi_n)_n$  of  $\mathbb{Q}\mathcal{L}$  such that  $\lim_n \varphi_n = \varphi$ .

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$\mathbb{R}\mathcal{L}^*$  is the logic obtained from  $\mathbb{R}\mathcal{L}$  by adding deduction rules

$$(\star) \quad \text{if } \varphi = \lim_m \varphi_m \quad \text{then} \quad \frac{\varphi_1, \varphi_2, \dots, \varphi_m, \dots}{\varphi}$$

## Limits and deduction

Let  $\varphi$  be a formula of  $\mathbb{R}\mathcal{L}$ . There exists a sequence of formulas  $\Theta = \{\varphi_n\}_{n \in \mathbb{N}} \subseteq \text{Form}_{\mathbb{Q}\mathcal{L}}$  such that:

- $\lim_n \varphi_n = \varphi$ ,
- $\text{Thm}(\varphi, \mathbb{R}\mathcal{L}^*) = \text{Thm}(\Theta, \mathbb{R}\mathcal{L}^*)$ .

## Theorem. TFAE:

- $\lim_n \varphi_n = \varphi$ ,
- $\lim_n f_{\varphi_n} = f_{\varphi}$  (uniform convergence),
- there exists an increasing sequence  $(r_n)_n$  in  $[0, 1]$  such that  $\bigvee_n r_n = 1$  and  $\vdash \eta_{r_n} \rightarrow (\varphi \leftrightarrow \varphi_n)$  for any  $n$

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How to axiomatize  $C([0, 1]^n)$ ?



### Marra V. and Reggio L., 2016

- They introduce  $\delta$ -algebra: MV-algebra endowed with an infinitary operator  $\delta$  such that, for sequences in  $C(X)$ , is  $\sum_n \frac{f_n}{2^n}$ .
  - The category of compact Hausdorff spaces is dual to the category of  $\delta$ -algebras with  $\delta$ -preserving MV-algebra morphisms.
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- Hence,  $\delta$ -algebras are equivalent to norm-complete Riesz MV-algebras,
  - In  $\mathbf{RL}$  : given  $\{\varphi_n\}_n$  we define the sequence  $\sigma_1 = \Delta_{\frac{1}{2}}\varphi_1$ ,  
 $\sigma_2 = \Delta_{\frac{1}{2}}\varphi_1 \oplus \Delta_{\frac{1}{2^2}}\varphi_2, \dots$  and we set  $\delta(\varphi_1, \varphi_2, \dots) = \lim_n \sigma_n$ .

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What happens if we close to  $\bigvee_n$  and  $\bigwedge_n$ ?

# The infinitary logic $\mathcal{IRL}$

- Language:  $\{\rightarrow, \neg\} \cup \{\nabla_r\}_{r \in [0,1]} \cup \vee$

- Axioms: the ones of  $\mathcal{RL}$  and

(S1)  $\varphi_k \rightarrow \bigvee_{n \in \mathbb{N}} \varphi_n$ , for any  $k \in \mathbb{N}$

- Deduction rules: Modus Ponens and

(SUP) 
$$\frac{(\varphi_1 \rightarrow \psi), \dots, (\varphi_k \rightarrow \psi) \dots}{\bigvee_{n \in \mathbb{N}} \varphi_n \rightarrow \psi}$$

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$$(SUP) \quad \frac{(\varphi_1 \rightarrow \psi), \dots, (\varphi_k \rightarrow \psi) \dots}{\bigvee_{n \in \mathbb{N}} \varphi_n \rightarrow \psi}$$

Following [Karp C. R., \*Languages with expressions of infinite length\*, North-Holland Pub. Co., 1964](#)

## The infinitary logic $\mathcal{IRL}$

- Models are  $\sigma$ -complete Riesz MV-algebras
- $\mathcal{IRL}$  is complete wrt to all models
- $Lind_{\mathcal{IRL}}$  is the Dedekind  $\sigma$ -completion of  $Lind_{\mathcal{RL}}$



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### A logic for quasi-Stonean spaces

Since any  $\sigma$ -complete Riesz MV-algebra is norm-complete (w.r.t. unit norm), any model of  $\mathcal{IRL}$  has the form  $C(X)$  where  $X$  is quasi-Stonean (basically disconnected compact Hausdorff space). In particular,  $Lind_{\mathcal{IRL}}$  has this form.

## The infinitary logic $\mathcal{IRL}$

- Models are  $\sigma$ -complete Riesz MV-algebras
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A *Riesz tribe* over  $X$  is a Riesz MV-algebra of  $[0, 1]$ -valued functions over  $X$  that are closed under pointwise countable suprema.

# The Loomis-Sikorski theorem for Riesz MV-algebras

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## The Loomis-Sikorski Theorem

Let  $A \subseteq C(X)$  be a  $\sigma$ -complete Riesz MV-algebra, where  $X = \text{Max}(A)$ . Then there exist a Riesz tribe  $\mathcal{T}$  and a  $\sigma$ -homomorphism of  $\mathcal{T}$  onto  $A$ .

## Representation using tribes

$\text{Lind}_{\text{IRL},n}$  is isomorphic with the Riesz tribe on  $[0, 1]^n$  generated by the projection functions.

## The infinitary logic $\mathcal{IRL}$

Słomiński J., *The theory of abstract algebras with infinitary operations*,  
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### Infinitary variety

The class of Dedekind  $\sigma$ -complete Riesz MV-algebras is the infinitary variety  $HSP([0, 1])$ .

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## Standard completeness for $\mathcal{IRL}$

If  $\varphi$  is a formula of  $\mathcal{IRL}$ , then  $\vdash_{\mathcal{IRL}} \varphi$  if and only if  $e(\varphi) = 1$  for any evaluation  $e : FORM_{\mathcal{IRL}} \rightarrow [0, 1]$ .

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- Neural networks as formulas of  $\mathbb{R}\mathcal{L}$ .

### Multilayer perceptron

A multilayer perceptron with  $l$  hidden layers,  $n$  inputs and one output can be represented as a function  $F : [0, 1]^n \rightarrow [0, 1]$  such that

$$F(x_1, \dots, x_n) = \phi \left( \sum_{k=1}^{n^{(l)}} \omega_{ok}^l \phi \left( \sum_{j=1}^{n^{(l-1)}} \omega_{kj}^{l-1} \phi \left( \dots \left( \sum_{i=1}^n \omega_{li}^1 x_i + b_i \right) \dots \right) \right) \right),$$

where

- $\phi : \mathbb{R} \rightarrow [0, 1]$  is the activation function (monotone-nondecreasing continuous function)
- $\omega_{ok}^l$  is the synaptic weight from neuron  $k$  in the  $l$ -th hidden layer to the single output neuron  $o$ ,
- $\omega_{kj}^{l-1}$  is the synaptic weight from neuron  $j$  in the  $(l-1)$ -th hidden layer to neuron  $k$  in the  $l$ -th hidden layer, and so on for the other synaptic weights.

## Neural networks as formulas of $\mathbb{RL}$

Let  $\mathcal{N}_n$  the class of multilayer perceptrons where the activation function is the piecewise linear function  $\varrho(x) = \max(\min(1, x), 0)$ , and the synaptic weights are real numbers.

**Theorem:**  $\mathcal{N}_n = \text{Free}_{RMV}(n)$ .

(a) For every  $l, n, n^{(2)}, \dots, n^{(l)} \in \mathbb{N}$ , and  $\omega_{ij}, b_i \in \mathbb{R}$ , the function  $F : [0, 1]^n \mapsto [0, 1]$  defined as

$$F(x_1, \dots, x_n) = \varrho \left( \sum_{k=1}^{n^{(l)}} \omega_{ok} \varrho \left( \sum_{j=1}^{n^{(l-1)}} \omega_{kj} \varrho \left( \dots \left( \sum_{i=1}^n \omega_{li} x_i + b_i \right) \dots \right) \right) \right),$$
 belongs to  $\text{Free}_{RMV}(n)$ .

(b) For any  $f \in \text{Free}_{RMV}(n)$ , there exist  $l, n, n^{(2)}, \dots, n^{(l)} \in \mathbb{N}$ , and  $\omega_{ij}, b_i \in \mathbb{R}$  such that

$$f(x_1, \dots, x_n) = \varrho \left( \sum_{k=1}^{n^{(l)}} \omega_{ok} \varrho \left( \sum_{j=1}^{n^{(l-1)}} \omega_{kj} \varrho \left( \dots \left( \sum_{i=1}^n \omega_{li} x_i + b_i \right) \dots \right) \right) \right).$$

Di Nola A., B. Gerla, I.L., 2013

### Problem

For  $f : [0, 1]^n \rightarrow \mathbb{R}$  an affine function, find a formula  $\varphi$  of  $\mathbb{RL}$  such that  $f_\varphi = \varrho \circ f$ .

```
function Formula( $(r_1, i_1), \dots, (r_m, i_m)$ )
{
(F1) if  $r_k \leq 0$  for any  $k \in \{1, \dots, m\}$  then return( $\perp$ );
(F2) find  $k \in \{1, \dots, m\}$  such that  $r_k > 0$ ;
      if  $i_k = 0$  then  $\psi := \mathbf{r}_k \top$  else  $\psi := \mathbf{r}_k x_{i_k}$ ;
(F3) if  $m = 1$  then return( $\psi$ );
(F4)  $\varphi = \text{Formula}((r_1, i_1), \dots, (r_{k-1}, i_{k-1}), (r_{k+1}, i_{k+1}), \dots, (r_m, i_m))$  ;
       $\chi = \text{Formula}((-r_1, i_1), \dots, (-r_{k-1}, i_{k-1}), (-r_{k+1}, i_{k+1}), \dots, (-r_m, i_m))$ 
      return(( $\varphi \oplus \psi$ )  $\odot \neg \chi$ )
}
```

We illustrate how the algorithm works on a simple example.

For  $f : [0, 1]^2 \rightarrow [0, 1]$ ,  $f(x_1, x_2) = x_2 - 0.3x_1$  we call the function  
function Formula((1, 2), (-0.3, 1))

{

(F2)  $k = 1$ ,  $r_k = 1$ ,  $i_k = 2$ ;  $\psi := \mathbf{1}x_2$ ;

(F4)  $\varphi = \text{Formula}((-0.3, 1))$  ;  $\chi = \text{Formula}((0.3, 1))$  ;

**return**(( $\varphi \oplus \psi$ )  $\odot \neg\chi$ )

}

Denote  $\perp = \neg(\varphi \rightarrow \varphi)$ . One can easily see that  $\varphi = \perp$  and  $\chi = \mathbf{0.3}x_1$ , so  
the function returns

$$(\varphi \oplus \psi) \odot \neg\chi = (\perp \oplus \mathbf{1}x_2) \odot \neg\mathbf{0.3}x_1$$

which is logically equivalent with  $x_2 \odot \neg\mathbf{0.3}x_1$ .

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